Implementation of CMOS Adder for Area & Energy Efficient Arithmetic Applications

Prachi B. Deotale*, Chetan G. Thote**

* Dept.of Electronics & telecommunication Engg., Dr.Bhausaheb Nandurkar College of Engg & Tech., Yavatmal, Maharashtra, India.
** (M.E Digital Electronics), Head of Dept.of Electronics & telecommunication Engg., Dr.Bhausaheb Nandurkar College of Engg & Tech., Yavatmal, Maharashtra, India.

Abstract- The most fundamental arithmetic operation is addition which is used in a digital data path logic system. Arithmetic and logic units, Microprocessors, etc. are some examples where we need to use arithmetic operations for processing data, for calculating addresses respectively. There are different architectures for building adder circuit. For example: 1) carry look ahead adder (CLA), 2) carry propagate adder (CPA), 3) carry save adder (CSA), & 4) carry select adder (CSLA). Among these different architectures CSLA is a particular way of implementing adder that performs addition rapidly and are used for faster addition in many data processing processors. From observation of the carry select adder architecture we can see that there is scope for modification in order to significantly minimize the area and power consumed by the circuit. In this work we are going to propose simple and efficient modification at gate-level structure in CSLA. The prime factor which hinders the faster operation of adder is the time taken in propagation of carry and this can be alleviated by using “Carry select adder”. Because carry select adder generates multiple carries & then select carry to generate the SUM [8]. CSLA is faster but it is not area efficient as multiple ripple carry adder (RCA) pairs are required to generate partial SUM & CARRY by considering carry input Cin 1 & Cin 0, after that multiplexer is used to select final outputs i.e SUM & CARRY. Basic idea of this work is to use binary to excess-1 converter (BEC) in place of RCA with input Cin=1 in the regular CSLA structure, so that lower area and power consumption can be achieved [2]-[4]. The prime advantage of using this BEC logic at the place of RCA with Cin 1 comes from the use of lesser number of logic gates than the number of logic gates required in n-bit full adder structure. The more elaboration about BEC logic is mentioned in following sections. For the purpose of explanation, we have used Square root carry select adder (SQRT CSLA) [5]-[6] in this paper.

II. METHODOLOGY FOR AREA & DELAY CALCULATION

There are three basic building blocks of CSLA

A) RCA
B) Multiplexer
C) BEC-1

A. XOR GATE:

In addition is the most fundamental arithmetic operation. Adders are the electronics circuits that perform the addition of numbers. In many data path logic systems, computers and processors adders are used for ALU, address calculation, increment and decrement operators, table indices calculation and for implementation of other arithmetic operations such as subtraction, multiplication and division etc. With the increase in chip density for implementation of more and more logical functions on a single chip the problem of area and power consumption is becoming more serious & it is the most dealt one by designers. For efficient & faster operation of VLSI systems there is a lot of research on design techniques is going on [1]. In this work we have proposed gate level modification in the architecture of carry select adder (CSLA). The prime factor which hinders the faster operation of adder is the time taken in propagation of carry and this can be alleviated by using “Carry select adder”. Because carry select adder generates multiple carries & then select carry to generate the SUM [8]. CSLA is faster but it is not area efficient as multiple ripple carry adder
Fig. 1 shows the implementation of AND, OR, INVERTER (AOI) implementation of XOR gate. The gates shown under dotted lines perform their operation in parallel. The numeric representation of each gate indicates the delay incorporated by that logic gate.

In this method of calculation we have considered the following points:
1. All gates used in implementation are AND, OR, INVERTER.
2. Each gate has gate delay equal to 1 unit.
3. We calculate delay by counting number of logic gates in the longest path of logic block & this delay contributes to maximum delay.
4. Area calculation is carried out by counting number of AOI gates required for implementation of each logic block.

Based on this methodology areas & delays of (2:1) MUX, XOR, HA & FA are calculated and results are tabulated in following table 1.

- **AREA**: Total number of AOI required implementing XOR logic block are 2INV, 2AND, 1OR gates. Hence (2+2+1=5) area count is 5.
- **DELAY**: Longest logic path in XOR logic block consist of 1INV, 1AND, & 1OR gate. Hence (1+1+1=3) delay is 3.

**B. BINARY TO EXCESS-1 CONVERTER**:

Fig 2. Shows the implementation of BEC-1 using XOR, AND & INV gates. The operation of the 4 bi BEC-1 can be represented by using function equations (1)-(4) given as below. (LEGENDS: ~NOT, &AND, ^OR)

\[
\begin{align*}
X_0 & = \neg B_0 \\
X_1 & = B_0 \oplus B_1 \\
X_2 & = B_2 \oplus (B_0 \& B_1) \\
X_3 & = B_3 \oplus (B_0 \& B_1 \& B_2)
\end{align*}
\]

- **AREA**: There are total 12 AOI gates are required for implementation of 4 bit BEC-1 (2AND, 3XOR (2INV, 2OR, 2AND), 1INV) (2*1+3*5+1*1=18). Hence area count is 18
- **DELAY**: There are 4 gates in the longest path of BEC-1 logic block (1XOR, 1AND) (1*3+1*1=4). Hence delay is 4.

Similarly using above method the area & delay count for different logic blocks that are used in a carry select adder structure are calculated & values are noted in Table 1.below.

<table>
<thead>
<tr>
<th>LOGIC BLOCK</th>
<th>DELAY</th>
<th>AREA</th>
</tr>
</thead>
<tbody>
<tr>
<td>XOR</td>
<td>3</td>
<td>5</td>
</tr>
<tr>
<td>2:1MUX</td>
<td>3</td>
<td>4</td>
</tr>
<tr>
<td>HA</td>
<td>3</td>
<td>6</td>
</tr>
<tr>
<td>FA</td>
<td>6</td>
<td>13</td>
</tr>
<tr>
<td>BEC1(4 BIT)</td>
<td>4</td>
<td>12</td>
</tr>
</tbody>
</table>

Based on the values derived in the previous section we can obtain the values for area count and area count of different groups of 32-bit CSLA block.

**III. 32-BIT REGULAR SQRT CSLA BLOCK**

**A. Structure of 32 bit regular SQRT CSLA**

The CSLA is used in many digital system designs to overcome the problem of carry propagation delay by
independently performing addition operation by considering carry inputs (Cin) as 1 and 0.

Fig. 3 shows a 32-bit Regular SQRT CSLA. The SQRT CSLA is divided into \( m=\sqrt{2m} \) carry select stages (CSS), where \( m \) is number of input bits. The 32 bit SQRT CSLA consists of 7 CSS. The CSS consists of two ripple carry adders one with carry in 0 and other with carry in 1. It also consists of a multiplexer which is used to select the sum and carry values from the two RCAs by using the control signal to it. The control signal to multiplexer is nothing but the carry out of the previous CSS. If the control signal is 1 then sum and carry out of RCA with Cin=1 is selected by the multiplexer and if control signal is 0 then sum and carry out of RCA with Cin=0 is selected by the multiplexer.

B. Calculation of area count for regular 32 bit SQRT CSLA

- Group 1

First group of 32 bit regular CSLA consist of 1 2-bit RCA. Fig. 4 Shows internal Structure of 2-bit RCA. It requires 1FA & 1HA. The area count of HA is 6 and that of FA is 13. Therefore the total area count for 2-bit RCA is 19.

As shown in fig. 5 Second group of CSLA consists of two 2-bit RCAs (for cin=0 & cin=1) & 1(6:3) MUX. One set of RCA for cin=1 consists of 2FA and other set of RCA for cin=0 consists of 1FA and 1HA. Based on area calculation method we obtain area count for group 2 is 57.

- Group 2

As shown in fig. 6. The third group of CSLA consists of two 3-bit RCAs (for cin=0 & cin=1) & 1(8:4) MUX. One set of RCA for cin=0 consists of 1HA & 2FA and other RAC for cin=1 consists of 3FA. The area of Group 3 is 87.

- Group 3

As shown in fig. 7. The fourth group of CSLA consists of two 4-bit RCAs (for cin=0 & cin=1) & 1(10:5) MUX. One set of
RCA for cin=1 consists of 4FA and other set of RCA for cin=0 consists of 1HA and 3HA. Hence area count for Group 4 is 117.

**Figure 7. Group 4 of 32-bit regular Sqrt CSLA**

Group 4 = 1HA + 7FA + 1(10:5) MUX

1HA = 6(1*6)
7FA = 91(7*13)
1(10:5) MUX = 20(5*4)

AREA = 117 (6 + 91 + 20)

Similarly, area count for all 8 groups of 32 bit Regular CSLA is calculated and the resulting values are mentioned in the table 2. Due to the use of two RCAs there is large requirement of the area, since RCA with cin=1 requires ‘n’ number of FAs for n-bit addition. Because of this reason, regular CSLA is faster but it is not area efficient. This lack of area efficiency can be alleviated by using BEC-1 at the place of RCA with cin=1.

IV. 32-BIT MODIFIED SQR T CSLA BLOCK

A. Structure of 32 bit modified SQRT CSLA

The fig 8. Above shows the structure of 32-bit Modified Sqrt CSLA. Comparing this modified structure with regular CSLA structure we can see that RCA with cin=1 is replaced by BEC-1 in the modified structure. For replacing n-bit RCA we require n+1 bit BEC-1 logic block. Let us see how it affects the values of area count for the different groups of Modified 32 bit CSLA structure.

**Figure 8. 32 bit Modified Sqrt CSLA block diagram.**

B. Calculation of area count for modified 32 bit SQRT CSLA

- Group 2

As shown in fig 9. Group 2 consists of 1 2-bit RCA with cin=0, 1 3-bit BEC-1, and 1 (6:3) MUX. The structure of 3-bit BEC-1 is built by using 2XOR gates, 1AND and 1NOT gate. The area count for group 2 is 43 as calculated below.

**Figure 9. Group 2 of 32-bit modified Sqrt CSLA**

Group 2 = 1FA + 1HA + 2XOR + 1AND + 1NOT + 1(6:3) MUX

1FA = 13(1*13)
1HA = 6(1*6)
2XOR = 10(2*5)
1AND = 1(1*1)
1NOT = 1(1*1)
1(6:3) MUX = 12(3*4)

AREA = 43 (13 + 6 + 10 + 1 + 1 + 12).

2) Group 3

As shown in fig 10. Group 3 consists of 1 (3-bit) RCA with cin=0, 1 (4-bit) BEC-1, 1(8:4) MUX. The structure of 4-bit BEC-1 is built by using 3XOR, 2AND & 1NOT gates. The area count for group 3 is 66 as calculated below.

**Figure 10. Group 3 of 32-bit modified Sqrt CSLA**

Group 3 = 3FA + 1HA + 4XOR + 3AND + 1NOT + 1(8:4) MUX

3FA = 39(3*13)
1HA = 6(1*6)
4XOR = 20(4*5)
3AND = 3(3*1)
1NOT = 1(1*1)
1(8:4) MUX = 20(5*4)

AREA = 89 (39 + 6 + 20 + 3 + 1 + 20)

Similarly, area count for remaining groups of Modified 32-bit CSLA is calculated and the resulting values are noted in table 2 below. As mentioned earlier we have proposed the use of BEC-1 in modified structure of 32-bit CSLA at the place of RCA with cin=1 in the regular structure.
Table 2. Area count for different groups of 32-bit SQRT CSLA

<table>
<thead>
<tr>
<th>GROUPS</th>
<th>REGULAR STRUCTURE</th>
<th>MODIFIED STRUCTURE</th>
</tr>
</thead>
<tbody>
<tr>
<td>GROUP 1</td>
<td>19</td>
<td>19</td>
</tr>
<tr>
<td>GROUP 2</td>
<td>57</td>
<td>43</td>
</tr>
<tr>
<td>GROUP 3</td>
<td>87</td>
<td>66</td>
</tr>
<tr>
<td>GROUP 4</td>
<td>117</td>
<td>89</td>
</tr>
<tr>
<td>GROUP 5</td>
<td>147</td>
<td>112</td>
</tr>
<tr>
<td>GROUP 6</td>
<td>177</td>
<td>135</td>
</tr>
<tr>
<td>GROUP 7</td>
<td>207</td>
<td>158</td>
</tr>
<tr>
<td>GROUP 8</td>
<td>87</td>
<td>66</td>
</tr>
<tr>
<td>TOTAL</td>
<td>898</td>
<td>688</td>
</tr>
</tbody>
</table>

** Calculation of reduction in area count: 898-688=210. Percentage of reduction in area count= 210/898*100= 23.38%. 

On comparing area count for regular and modified structure of 32-bit SQRT CSLA, it is evident that the area count of CSLA is reduced by 210 i.e. around 23.38% of area occupancy could be reduced by using modified structure for 32-bit CSLA implementation.

V. SIMULATION RESULTS AND COMPARISON

For the simulation of 32-bit modified SQRT CSLA architecture we have used Tanner EDA Tool (v.13), design is carried out using 0.25µm technology. The simulation results for area and power of regular 32-bit CSLA are taken as a reference from previous research papers [7].

A. Results of simulation

Fig 11. Shows the resulting schematic of 32-bit Modified Sqrt CSLA with inputs A= FFFF_FFFF & B=FFFF_FFFF.

Fig 12. Shows schematic of the internal structure of 32-bit modified Sqrt CSLA. We can see different groups of CSLA containing RCAs, BEC-1s & MUXs. Here RCA with Cin=1 is replaced by BEC-1.

Table 3. Results in tabular form

<table>
<thead>
<tr>
<th>Parameters</th>
<th>Resulting values</th>
</tr>
</thead>
<tbody>
<tr>
<td>Total no. of devices</td>
<td>2356(MOSFETs)</td>
</tr>
<tr>
<td>Voltage</td>
<td>5 volts</td>
</tr>
<tr>
<td>Current</td>
<td>152.5492 nA</td>
</tr>
<tr>
<td>Power</td>
<td>762.746nW</td>
</tr>
<tr>
<td>Simulation delay</td>
<td>217.25secs ~3 mint.</td>
</tr>
<tr>
<td>Temp.</td>
<td>25° c.</td>
</tr>
</tbody>
</table>
B. Comparison

Table 4 shows the comparison of regular & modified 32-bit CSLA based on area and power consumption.

Table 4. Comparison table

<table>
<thead>
<tr>
<th>Parameters</th>
<th>Regular 32-bit CSLA</th>
<th>Modified 32-bit CSLA</th>
</tr>
</thead>
<tbody>
<tr>
<td>Area (gate count)</td>
<td>898</td>
<td>688</td>
</tr>
<tr>
<td>Power</td>
<td>1127.3µW</td>
<td>762.746nW</td>
</tr>
</tbody>
</table>

V. CONCLUSION

In this work we have used gate level modification i.e BEC-1 logic is used at the place of RCA with cin=1 in regular CSLA to create modified structure. The resulting implementation shows that the area consumption has reduced by 23% and power consumption is also reduced due to use of less number of gates (Results in Table 3&4). With this technique of adder implementation we can achieve more efficient structures for 64-, 128-bit adders which can give use significant reduction in area and power consumption & therefore more efficient arithmetic operation.

REFERENCES


AUTHORS

First Author – Prachi B. Deotale, Dept.of Electronics & telecommunication Engg., Dr.Bhausaheb Nandurkar College of Engg & Tech., Yavatmal, Maharashtra, India., prachi.deotale@gmail.com

Second Author – Chetan G. Thote, (M.E Digital Electronics), Head of Dept.of Electronics & telecommunication Engg., Dr.Bhausaheb Nandurkar College of Engg & Tech., Yavatmal, Maharashtra, India., chetan.thote86@gmail.com